



Stochastic Ordering for Internet Congestion Control

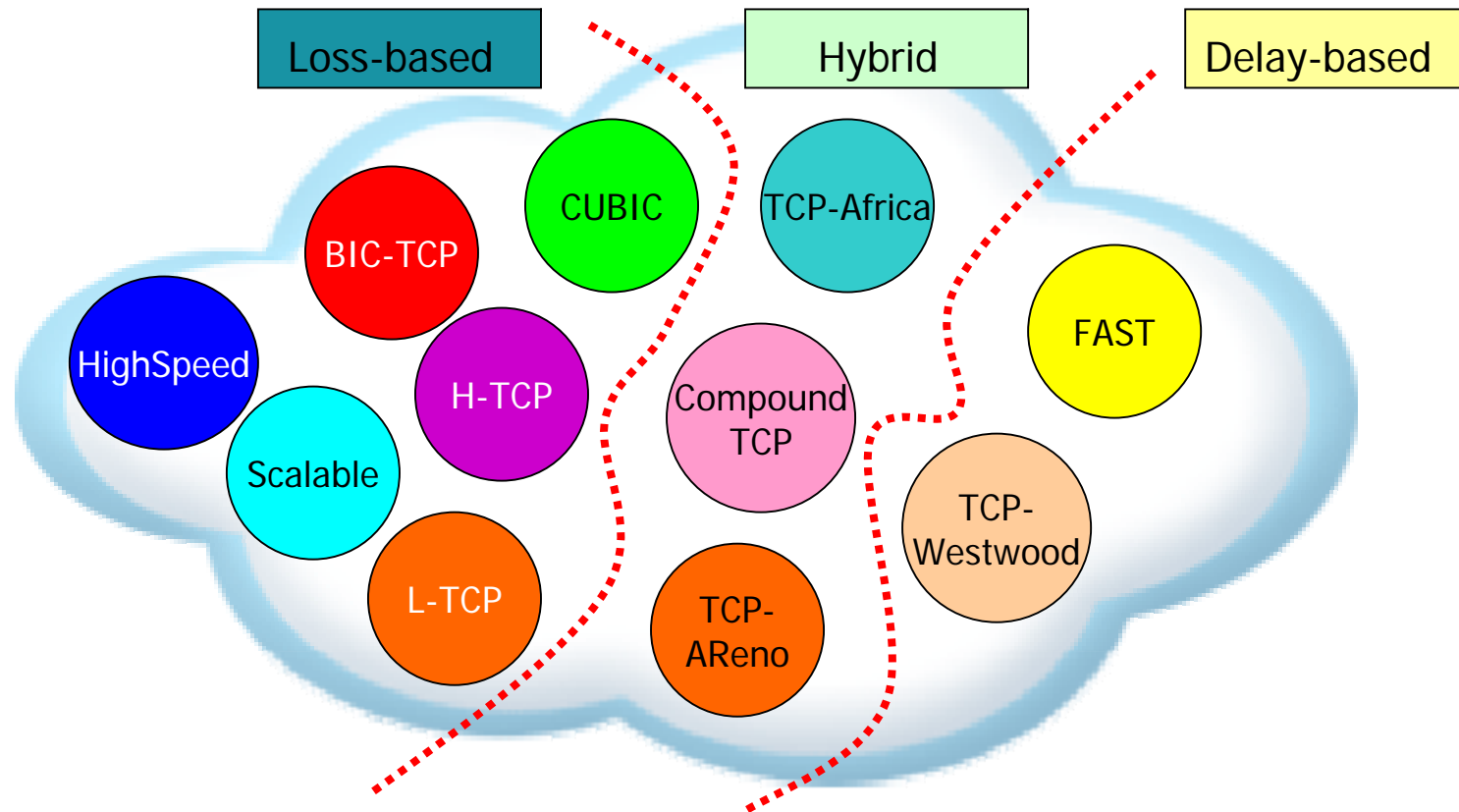
Han Cai, Do Young Eun, **Sangtae Ha**,
Injong Rhee, and Lisong Xu

PFLDnet 2007

February 7, 2007

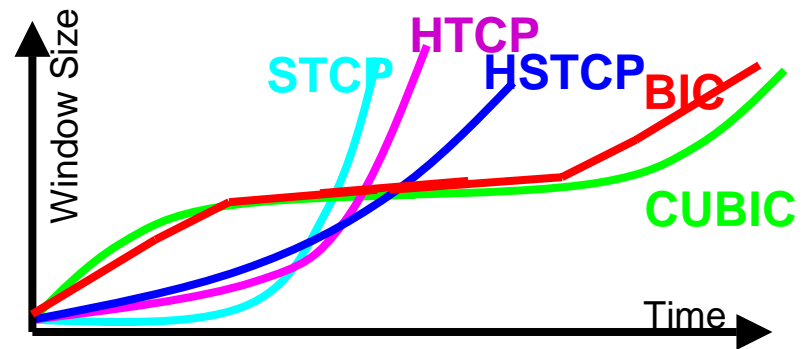
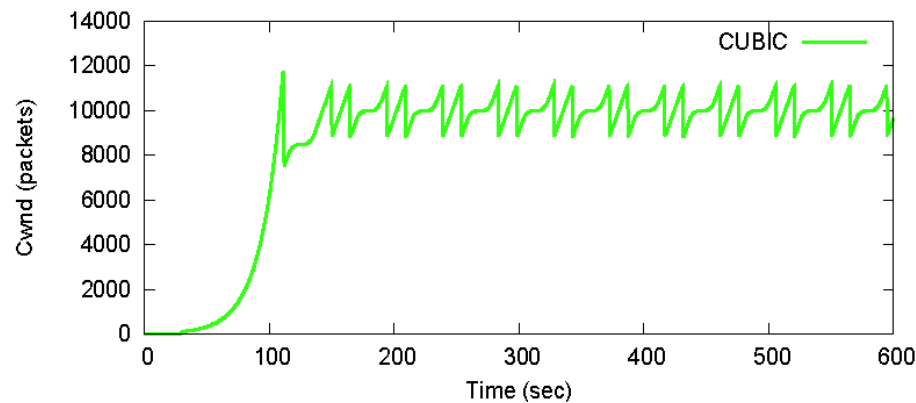
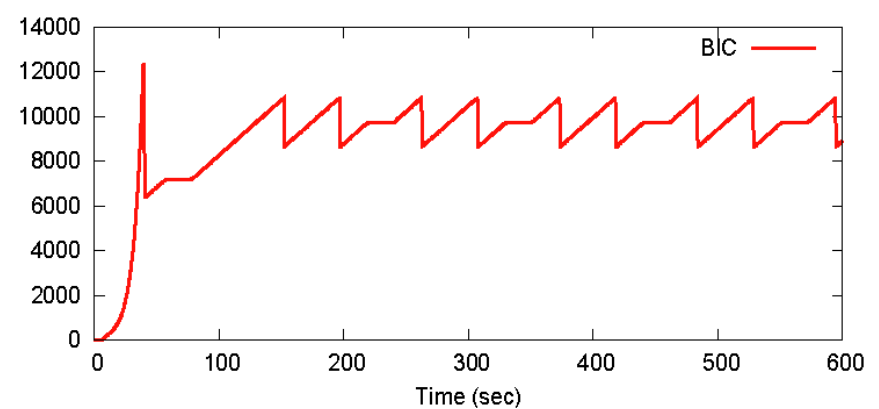
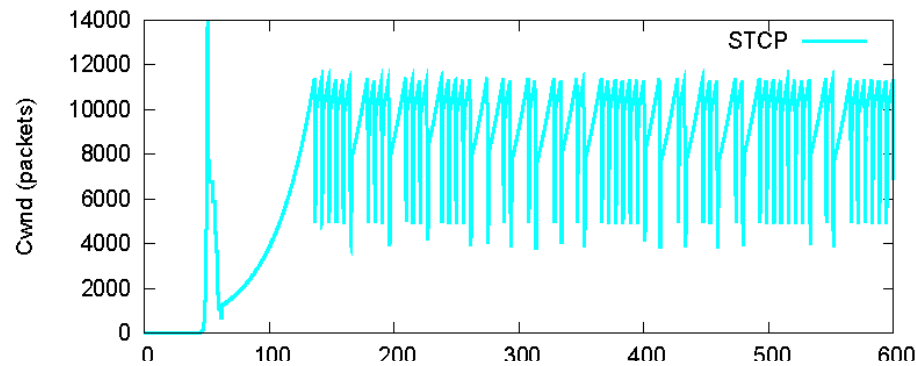
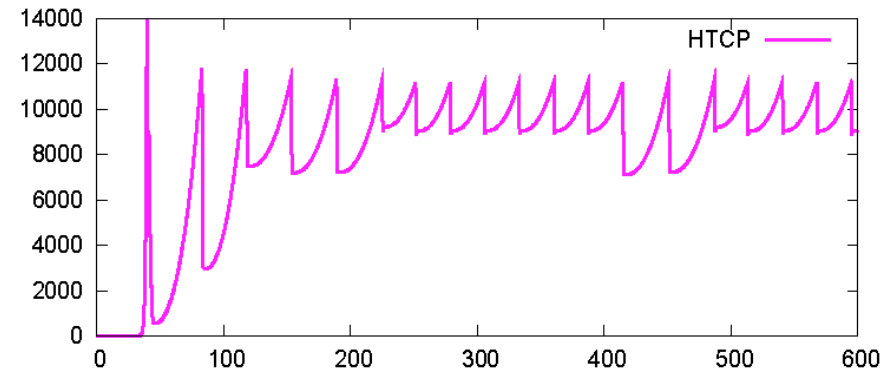
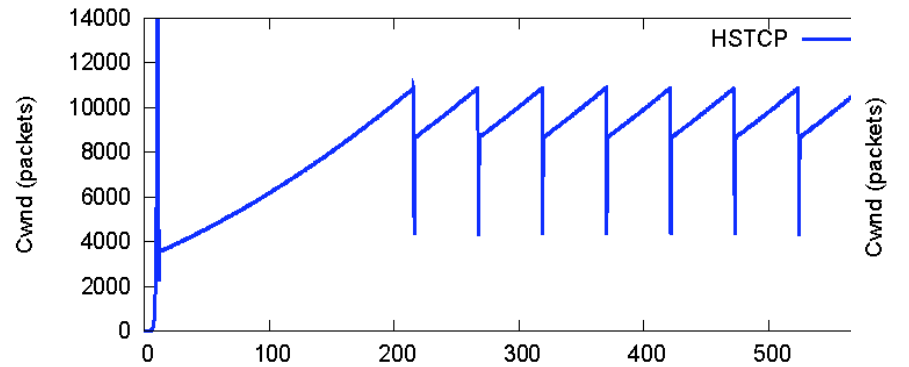


High-Speed TCP Variants



- Many High-speed TCP variants have been proposed
- High-speed protocols can be divided into three categories
- Can we compare loss-based protocols?

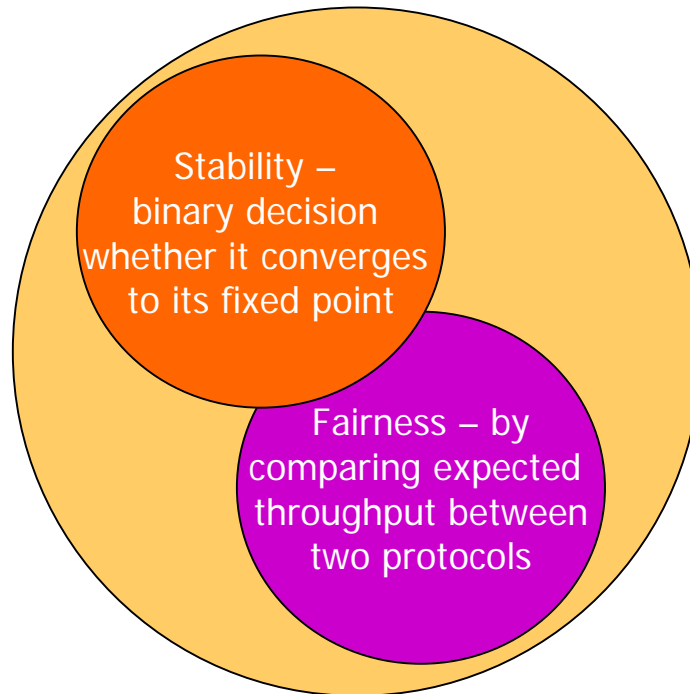
Example of Growth Functions



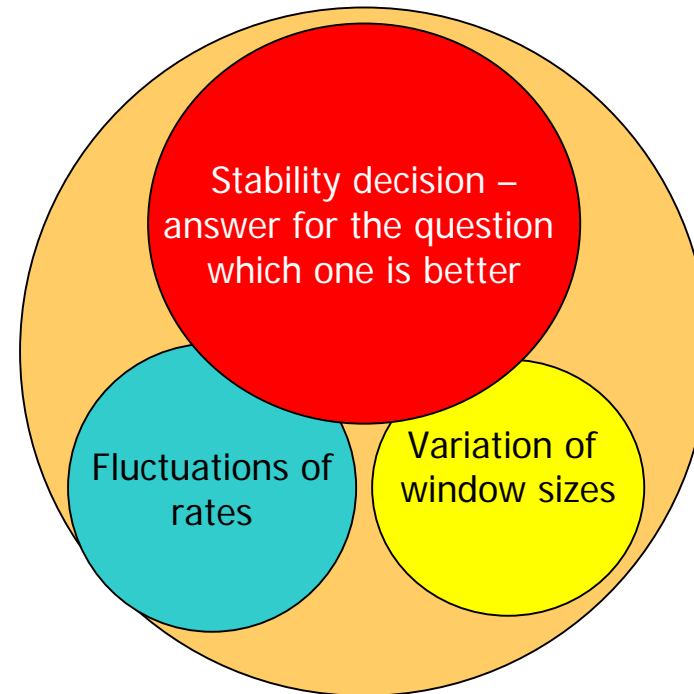


First Order vs. Second or High Order Behaviors

First order analysis

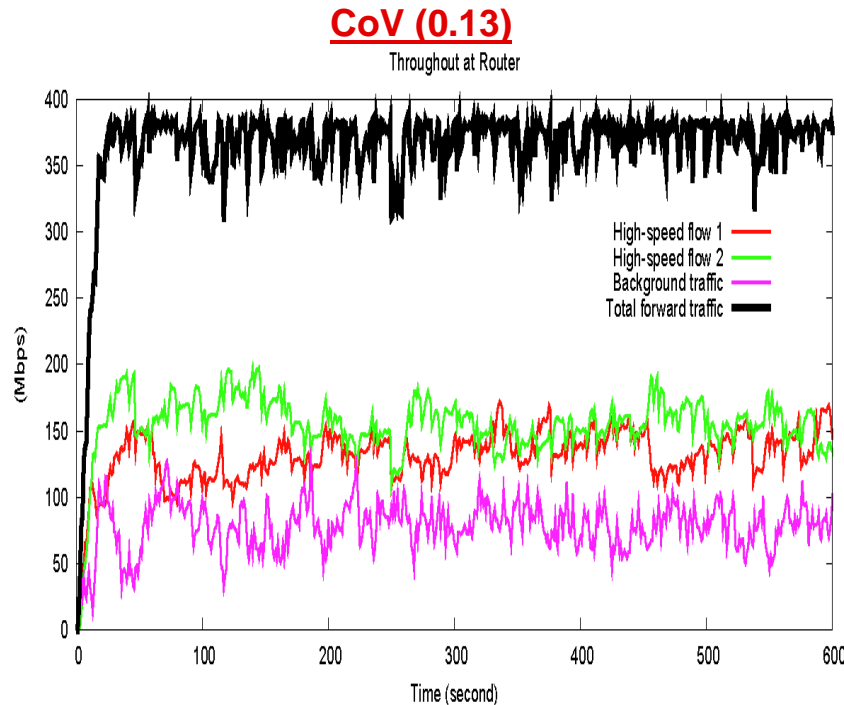


Second or high-order analysis

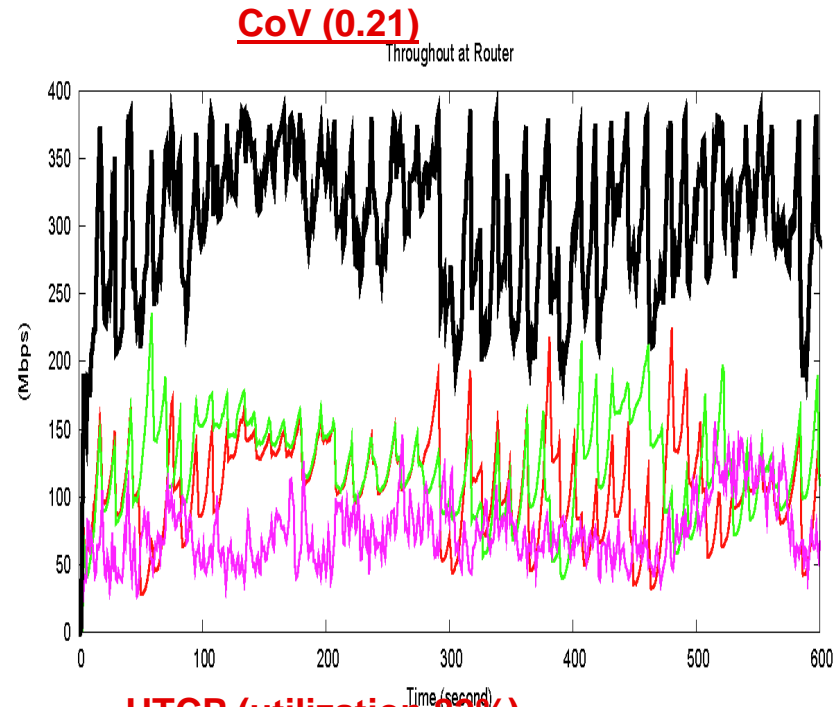


- If two protocols have the same average behaviors, they are not distinguishable by the first order analysis (fixed point analysis)
- Even if two protocols have the same first order behaviors (expected throughput), they could have widely different second or high-order behaviors.

CoV – one example of second order



BIC (utilization 94%)



HTCP (utilization 82%)

- CoV, standard deviation over mean, is commonly used by practitioners to compare the stability of protocols.
- Higher CoV also affects the general well beings of the network including utilization, queue oscillation, packet loss characteristics



Stochastic Convex Ordering

theoretical tool to analyze the 2nd order.

Definition: for random variables, X and Y,

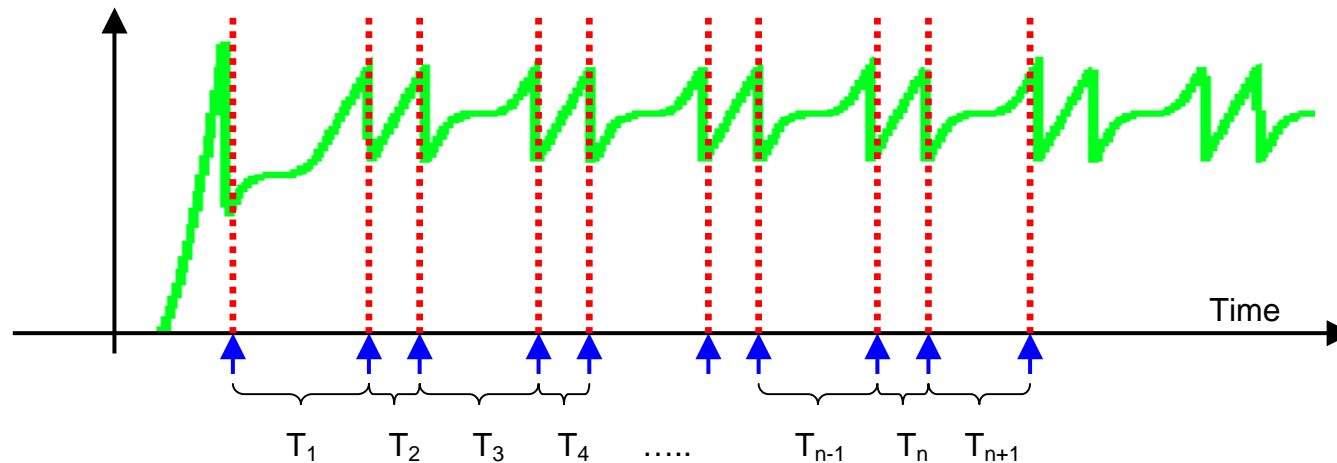
$$X \leq_{\text{cx}} Y \text{ means } E\{f(X)\} \leq E\{f(Y)\}$$

for any convex function f,
e.g, f = variance.

- If $X \leq_{\text{cx}} Y$ in convex ordering, it means X is less variable than Y.



Main Assumption



Assumption: The **Inter-loss intervals** T_n are **i.i.d** (independent and identically distributed) – commonly observed by many Internet measurement studies [4,5,6].

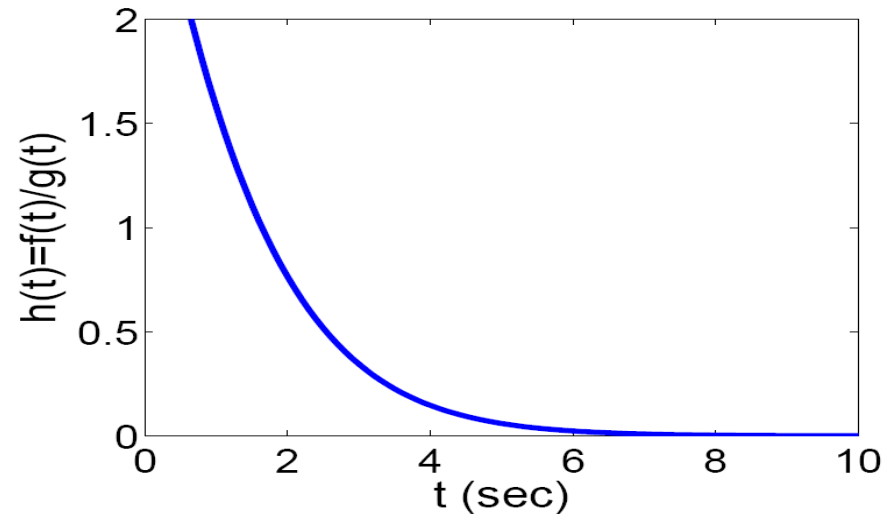
[4] Y. Zhang, N. Duffield, V. Paxson, and S. Shenker, "On the constancy of Internet path properties," ACM SIGCOMM IMW 2001

[5] E. Altman, K. Avrachenkov, and C. Barakat, "A Stochastic Model of TCP/IP with Stationary Random Loss," ACM SIGCOMM 2000

[6] J. Padhye, V. Firoiu, D. Towsley, and J. Kurose, "Modeling TCP Throughput: a Simple Model and its Empirical Validation," ACM SIGCOMM 1998



Main result



(a) monotone $h(t)$

- Given window profiles $f(t)$ and $g(t)$, $h(t) = f(t)/g(t)$ is monotonically decreasing,

then $f(t) \leq_{cx} g(t)$. Concave < Convex

- E.g.

$$f(t) = \log(t)$$

$$g(t) = t^2$$

$$f(t) = t^{0.5}$$

$$g(t) = t^2$$

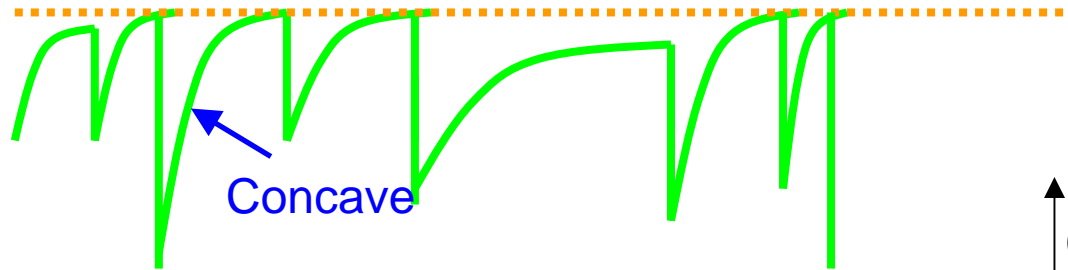
$$f(t) = 0.5t^2$$

$$g(t) = 3t^3$$

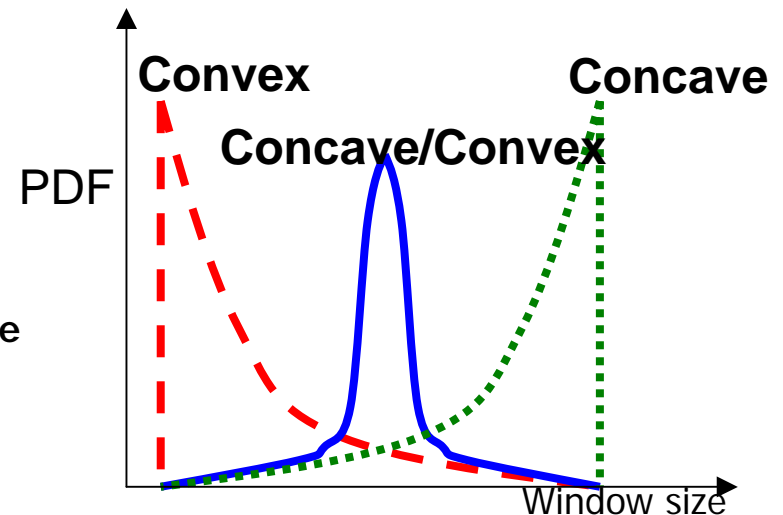
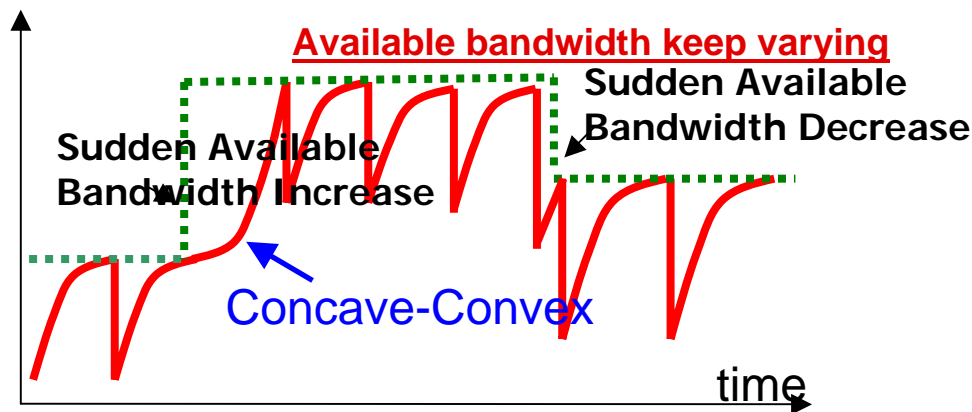


Results - Intuition

Stationary loss



Non-stationary loss



- During stationary loss, concave has less variance because its windows are mostly around the mean so that its variance is small.
- During non-stationary loss, concave-convex has also its windows mostly around the mean.

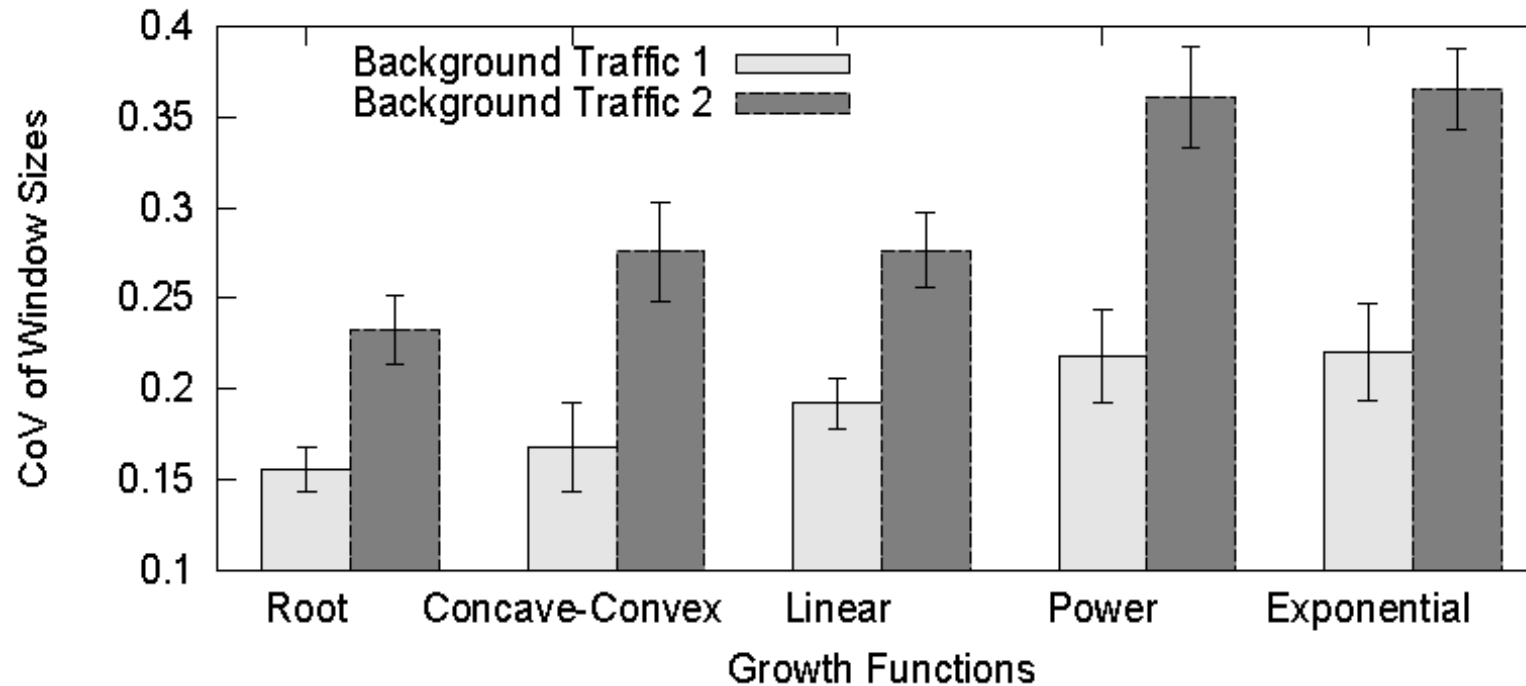


NS2 Simulation verification

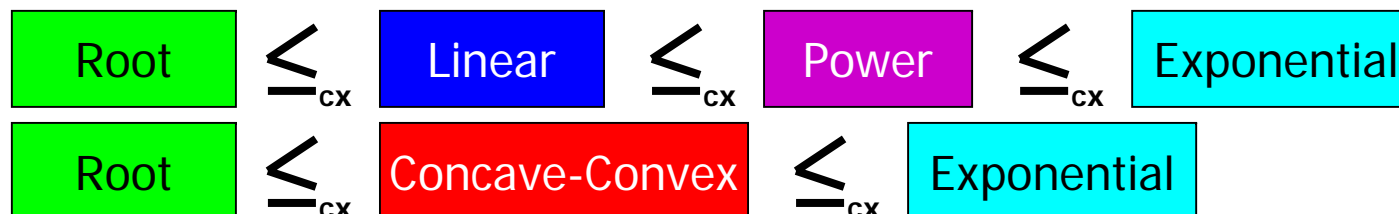
- Dumbbell, bottleneck 250Mbps, RTT 100ms, 100% BDP buffer size
- Loss generated by predefined models and by using background traffic
- Background Traffic (20% of total link bandwidth)
 - Type I (five long-lived flows), Type II (300 web sessions)
- Five pseudo protocols simulated
 - Root, Linear, Power (Square), Exponential, Concave-Convex
- Measure the CoVs of window sizes of the five pseudo protocols.



NS2 Simulation Result

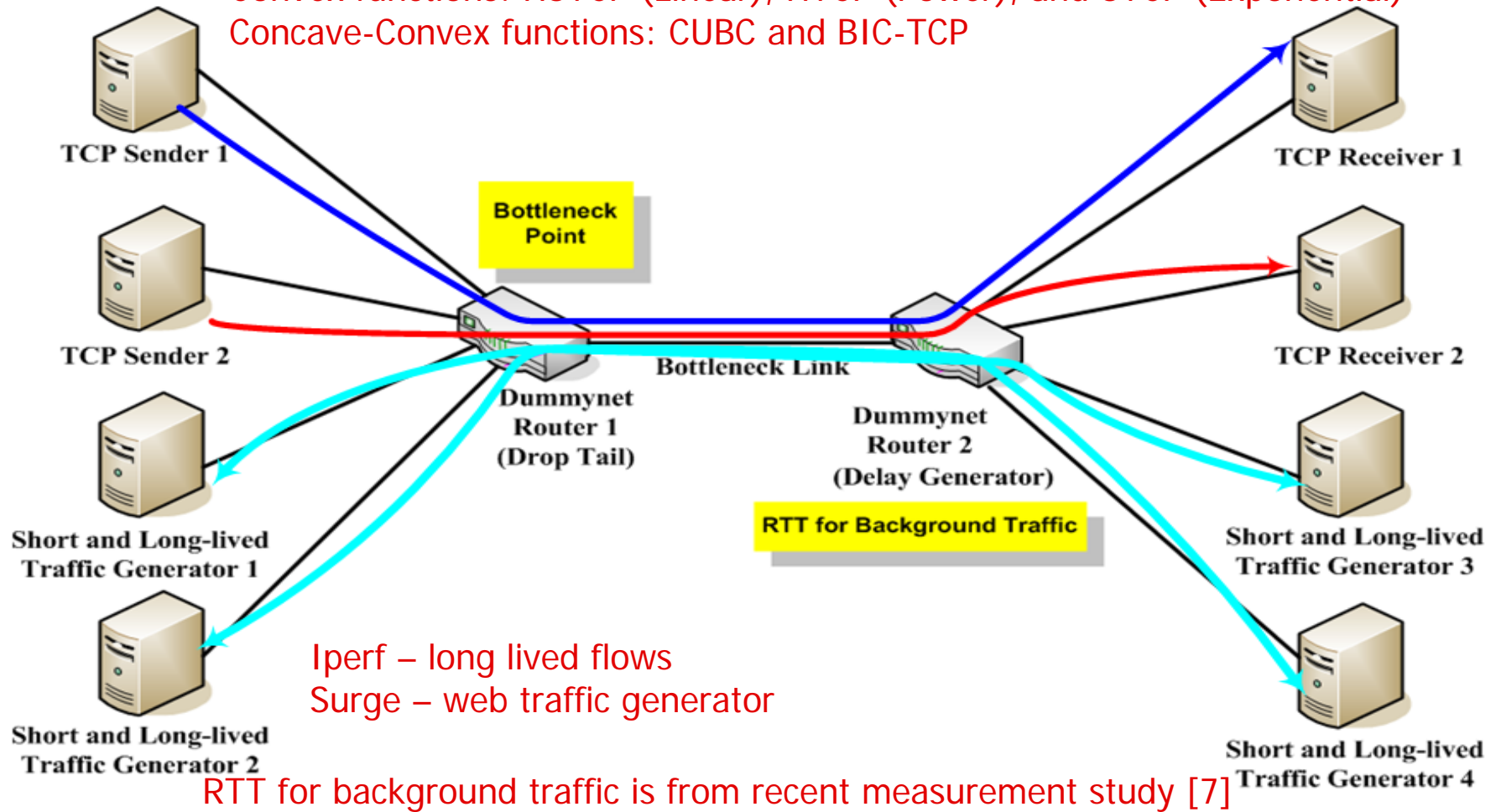


- The result confirms the same ordering predicted by our analytical result



Testbed (Dummynet) Setup

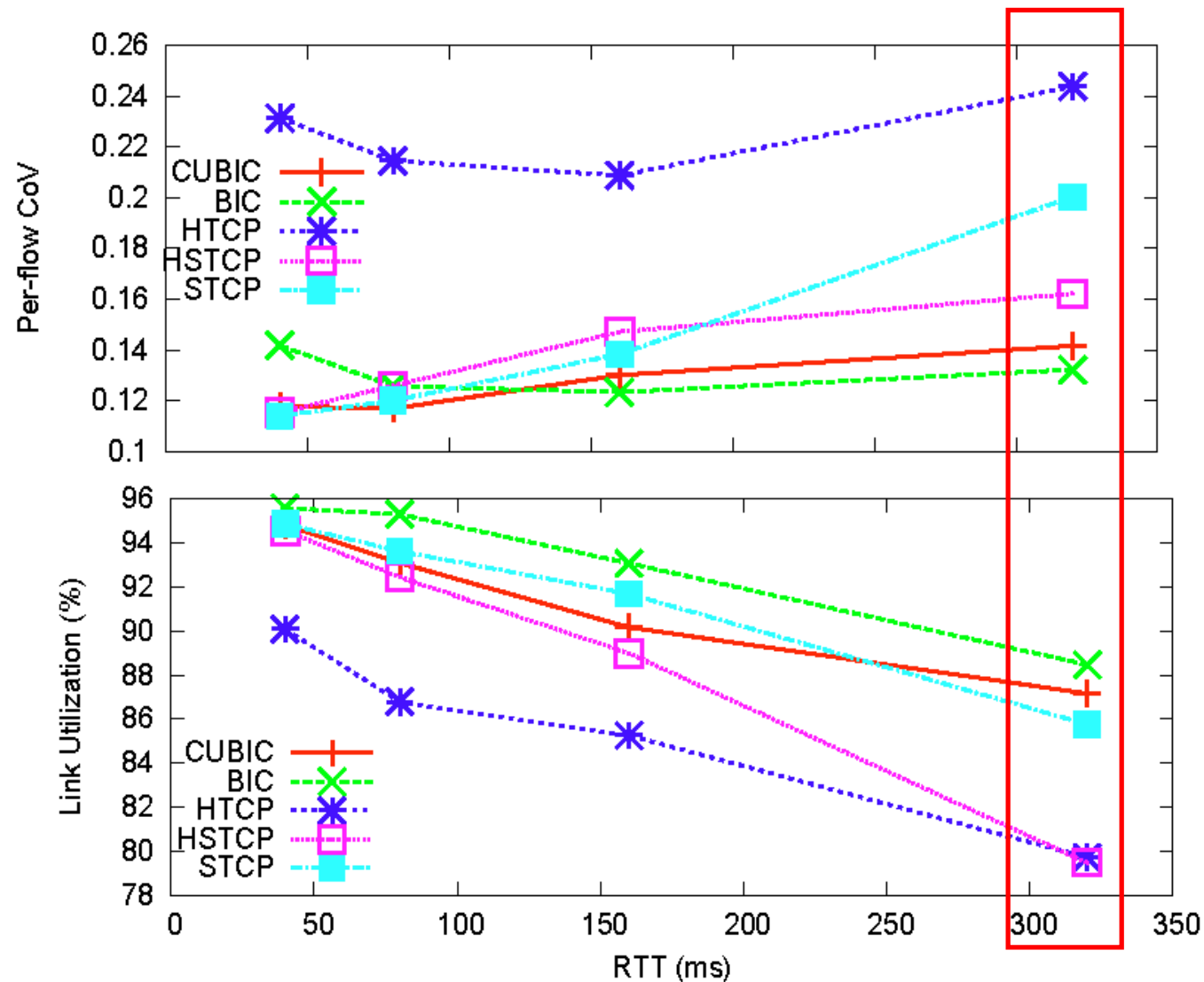
400Mbps Bottleneck,
 Convex functions: HSTCP (Linear), HTCP (Power), and STCP (Exponential)
 Concave-Convex functions: CUBC and BIC-TCP



[7] J. Aikat, J. Kaur, F. D. Smith, and K. Jeffay, Variability in TCP round-trip times, IMC 2003



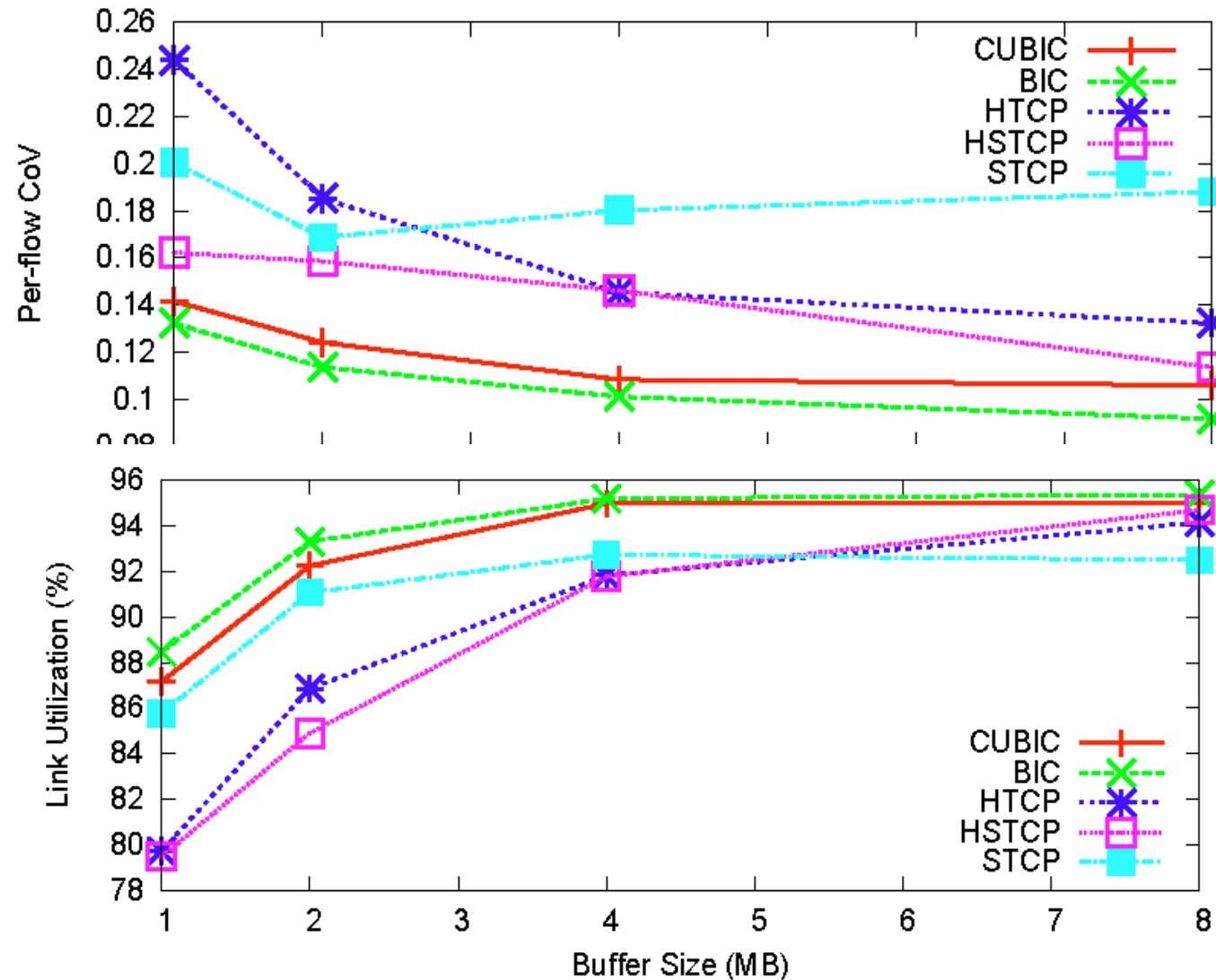
CoV and Link Utilization



- Buffer size (1MB), four HS flows with the same RTT (40ms – 320ms)



Impact of buffer sizes



- Buffer size (1 – 8MB), four HS flows with the same RTT (320ms)



Conclusion and Future Work

- Window growth function determines its relative stability.
- Stochastic Convex Ordering can be applicable to loss-based protocols.
- Concave-Convex protocols tends to give the smallest rate variation (BIC, CUBIC).
- Rate variations can affect the general well-beings of the network including utilization, queue oscillations and packet loss characteristics.
- Dynamics of aggregated flows and their impact on the general health of the networks would be our future work.



Q & A

More experimental results are available at
<http://netsrv.csc.ncsu.edu/convex-ordering>

Thank you for your participation