Optical XCP for Small Buffered Optical Packet Switching Networks

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Outline

- Problems in OPS networks
- Architecture
- Objective
- Proposed Solutions
- Simulation Results
- Conclusions

Problems in OPS Networks

Optical RAM is not available

- Electronic RAM is not a feasible solution
- O/E/O conversion is hard at ultra-high speed of optical networks
- Possible to use FDLs (fiber-delay-lines) for buffering
- FDLs can provide small amount of buffering with fixed delays

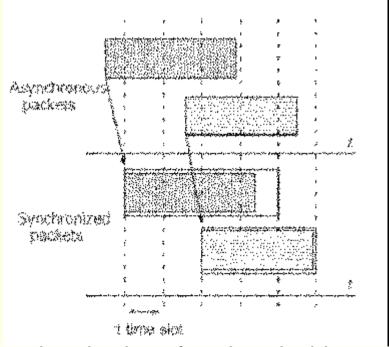
Switching and Processing

- Hard to use complex algorithms
- All-optical switching is necessary

Architecture

OPS Architecture

- Time-slotted WDM OPS network
- Variable length IP packets enter OPS domain without aggregation

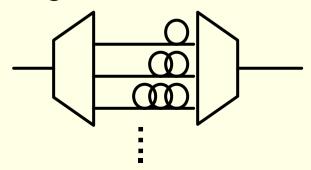


Synchronization of packets inside a router

Architecture

FDL Architecture

- Single stage FDL set with B delay lines
- FDL length distribution increases linearly (x,2x,3x,4x....) where x is FDL granularity
 - » FDL granularity of 3 means a FDL set of (3 slots, 6 slots, 9 slots, 12 slots....)
- Voids occur between packets inside FDLs when x>1
- No void-filling
- Output buffering



Objective

- Designing an all-optical OPS architecture that can achieve high utilization and low packet drop ratio by using small buffers
- Showing the buffer requirements

Advantages

- Decreasing the buffer requirements in the core
- Realizing all-optical high-speed OPS networks

Proposed Solutions 1/2

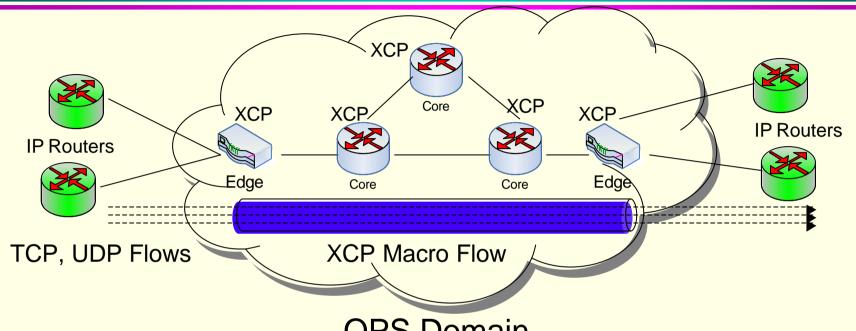
Preventing wavelength over-utilization

- Apply XCP-based congestion control
- Carefully select XCP parameters
- Control maximum wavelength utilization ratio

XCP feedbacks

- Carry feedbacks in probe packets
- Possible to use a control wavelength for feedbacks
- Control wavelength can be slow enough for O/E/O conversion, updating headers and calculating feedbacks

Proposed Solutions 2/2



OPS Domain

Burstiness

- Establish macro flows between edge nodes
- Assign incoming TCP, UDP flows to macro flows (similar to XCP-CSFQ, TeXCP)
- Apply token-based leaky bucket pacing to macro flows according to XCP flow rate at edge node
- Possible to use LSPs for controlling macro flows if GMPLS is available

Selecting Target Utilization 1/2

- At each XCP core node, wavelength capacity must be explicitly given to XCP control algorithm
- By giving a false capacity value to XCP lower than actual wavelength capacity:
 - XCP algorithm converges to the given false capacity
 - Possible to control maximum wavelength utilization
 - Can be explicitly given as a target utilization

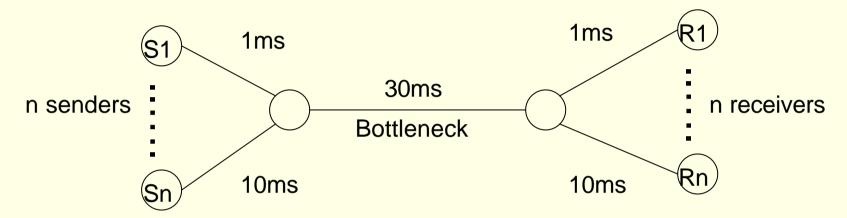
Selecting Target Utilization 2/2

Careful selection of target utilization can prevent

- Effective load overshoots of core routers due to voids in FDLs
- Utilization overshoots due to congestion and oscillations caused by XCP algorithm
- Core buffer buildups

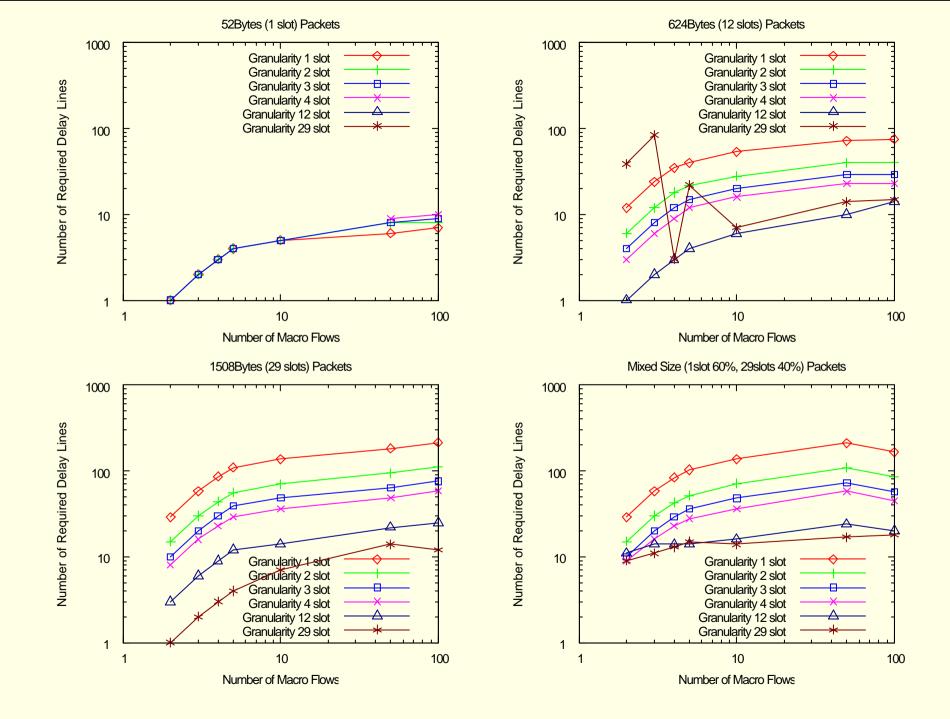
Simulation Settings

Slotted OPS with single data wavelength

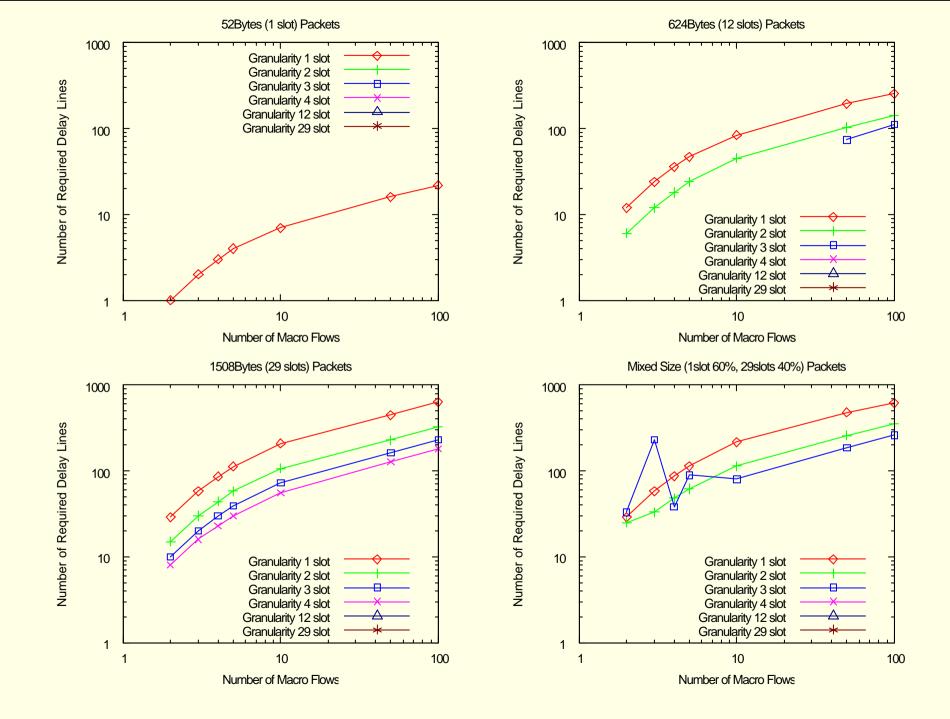


- Wavelength speed ranging from 172Mbit/s to 5Gbit/s
- Simulation duration is 40seconds
- Slot size is 52 Bytes
- Edge XCP agents always have packet to send
- Packet size distribution
 - All packets are 1 slot size (52 Bytes)
 - All packets are 29 slot size (1508 Bytes)
 - All packets are 12 slot size (624 Bytes)
 - Mixed Size (60% of packets are 1 slot size, 40% of packets are 29 slot size)

FDL Requirements (Target Utilization 30%)



FDL Requirements (Target Utilization 90%)



Conclusions and Future Work

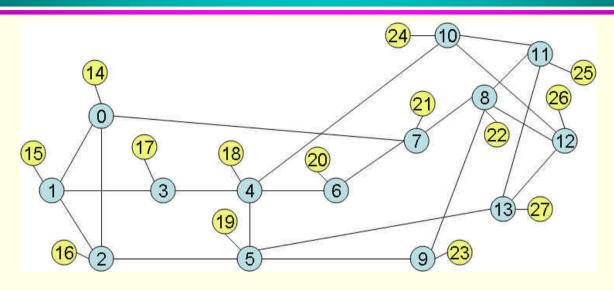
Conclusions

- Possible to control FDL requirements by applying pacing at edge, controlling the number of LSPs on a wavelength and limiting wavelength utilization
- XCP allows limiting utilization of wavelengths at a stable level in order to prevent queue-buildups
- Voids in FDLs can make the system unstable and cause drastic packet loss ratios
- Necessary to consider possible packet size distributions in order to minimize buffer requirements and guarantee stable operation

Future Work

- Multistage FDLs can further decrease the FDL requirements.
- Optimum size of electronic buffers on edge nodes

NSFNET Simulation Settings



- 28 nodes (14 edge + 14 core) and 35 links (21 core + 14 edge)
- Core nodes use only FDL for buffering
- Each edge source node sends traffic to all other edge nodes
- Edge XCP agents always have packet to send to other edge nodes
- Simulation duration is 40 seconds
- Edge XCP agents start randomly in the first 10seconds
- Wavelength speed ranging from 70Mbit/s to 2Gbit/s depending on simulated packet size distribution

